### Last Classes: Chart parsing and chunking

#### Today:

- 1. The Earley Algorithm
- 2. Probabilistic Parsing

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### **Efficient Parsing**

n =sentence length

Time complexity for naive algorithm: exponential in nTime complexity for bottom-up chart parser:  $\bigcirc(n^3)$ 

Options for improving efficiency:

1. Don't do twice what you can do once.

2. Don't represent distinctions that you don't need.

Fall leaves fall and spring leaves spring.

3. Don't do once what you can avoid altogether.

The can holds the water. ("can": AUX, V, N)  $\,$ 

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## Earley Algorithm: Top-Down Chart Parser

For all S rules of the form  $S \to X_1 \dots X_k$ , add a (top-down) edge from 1 to 1 labeled:  $S \to \circ X_1 \dots X_k$ .

Do until there is no input left:

- 1. If the agenda is empty, look up word categories for next word, add to agenda.
- 2. Select a constituent from the agenda: constituent C from  $p_1$  to  $p_2$ .
- 3. Using the (bottom-up) edge extension algorithm, combine C with every active edge on the chart (adding C to chart as well). Add any new constituents to the agenda.
- 4. For any active edges created in Step 3, add them to the chart using the top-down edge introduction algorithm.

Top-down edge introduction.

To add an edge  $S \to C_1 \dots \circ C_i \dots C_n$  ending at position j:

For each rule in the grammar of form  $C_i \to X_1 \dots X_k$ ,

recursively add the new edge  $C_i \to \circ X_1 \dots X_k$  from j to j.

Grammar Lexicon1. $S \rightarrow NP VP$ the: ART2. $NP \rightarrow ART ADJ N$ large: ADJ3. $NP \rightarrow ART N$ can: N, AUX, V4. $NP \rightarrow ADJ N$ hold: N, V5. $VP \rightarrow AUX VP$ water: N, V6. $VP \rightarrow V NP$ Sentence: $_1$ The $_2$ large $_3$ can $_4$ can $_5$ hold $_6$ water $_7$	<b>Probabilistic Parsing</b> Goal: Find the most likely parse. 1. Parsing with PCFGs 2. Problems 3. Probabilistic lexicalized CFGs
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<ul> <li>CFG's</li> <li>A context free grammar consists of: <ol> <li>a set of non-terminal symbols N</li> <li>a set of terminal symbols Σ (disjoint from N)</li> <li>a set of productions, P, each of the form A → α, where A is a non-terminal and α is a string of symbols from the infinite set of strings (Σ ∪ N)</li> <li>a designated start symbol S</li> </ol> </li> </ul>	Probabilistic CFGs         Augments each rule in P with a conditional probability: $A \rightarrow \beta [p]$ where p is the probability that the non-terminal A will be expanded to the sequence $\beta$ . Often referred to as $P(A \rightarrow \beta)$ or $P(A \rightarrow \beta   A)$ .

Example $S \rightarrow NP VP$ [.80] $Det \rightarrow that [.05]$ $the [.80]$ $a [.15]$ $S \rightarrow Aux NP VP$ [.15] $Noun \rightarrow book$ [.10] $S \rightarrow VP$ [.05] $Noun \rightarrow flights$ [.50] $NP \rightarrow Det Nom$ [.20] $Noun \rightarrow meal$ [.40] $NP \rightarrow Proper-Noun$ [.35] $Verb \rightarrow book$ [.30] $NP \rightarrow Pronoun$ [.40] $Verb \rightarrow want$ [.40] $Nom \rightarrow Noun$ [.55] $Aux \rightarrow can$ [.40] $Nom \rightarrow Noun Nom$ [.20] $Aux \rightarrow does$ [.30] $Nm \rightarrow Noun Nom$ [.55] $Proper-Noun \rightarrow TWA$ [.40] $VP \rightarrow Verb$ [.55] $Proper-Noun \rightarrow Denver$ [.40] $VP \rightarrow Verb NP$ [.40] $Proper-Noun \rightarrow Denver$ [.40] $VP \rightarrow Verb NP NP$ [.40] $Pronoun \rightarrow you[.40]$ $I[.60]$	<ul> <li>Why are PCFGs useful?</li> <li>Assigns a probability to each parse tree T</li> <li>Useful in disambiguation <ul> <li>Choose the most likely parse</li> <li>Computing the probability of a parse</li> <li>If we make independence assumptions, P(T) = ∏<sub>n∈T</sub> p(r(n)).</li> </ul> </li> <li>Useful in language modeling tasks</li> </ul>
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Example(a) $\overbrace{V}$ $\overbrace{V$ $\overbrace{V$ $\overbrace{V$ $\overbrace{V}$ $V$ $\overbrace{V$ $\overbrace{V$	Where do the probabilities come from? 1. from a treebank: $P(\alpha \rightarrow \beta   \alpha) = Count(\alpha \rightarrow \beta)/Count(\alpha)$ 2. use EM (forward-backward algorithm, inside-outside algorithm)
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## Parsing with PCFGs

Produce the most likely parse for a given sentence:

 $\hat{T}(S) = argmax_{T \in \tau(S)} P(T)$ 

where  $\tau(S)$  is the set of possible parse trees for S.

• Augment the Earley algorithm to compute the probability of each of its parses.

When adding an entry E of category C to the chart using rule i with n subconstituents,  $E_1, \ldots, E_n$ :

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 $P(E) = P(rule \ i \mid C) * P(E_1) * \dots * P(E_n)$ 

• probabilistic CYK (Cocke-Younger-Kasami) algorithm

## Problems with PCFGs

Do not model structural dependencies.

Often the choice of how a non-terminal expands depends on the location of the node in the parse tree.

E.g. Strong tendency in English for the syntactic subject of a spoken sentence to be a pronoun.

- 91% of declarative sentences in the Switchboard corpus are pronouns (vs. lexical).
- In contrast, 34% of direct objects in Switchboard are pronouns.

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Problems with PCFGs

Do not adequately model *lexical dependencies*.

Moscow sent more than 100,000 soldiers into Afghanistan...

PP can attach to either the NP or the VP: NP  $\rightarrow$  NP PP or VP  $\rightarrow$  V NP PP?

Attachment choice depends (in part) on the verb: *send* subcategorizes for a destination (e.g. expressed via a PP that begins with *into* or *to* or ...).