

Last Class: Parsing Intro

1. Grammars and parsing

Today: Parsing Algorithms

1. Top-down and bottom-up parsing
2. Chart parsers
3. Bottom-up chart parsing

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CFG's

A context free grammar consists of:

1. a set of non-terminal symbols N
2. a set of terminal symbols Σ (disjoint from N)
3. a set of productions, P , each of the form $A \rightarrow \alpha$, where A is a non-terminal and α is a string of symbols from the infinite set of strings $(\Sigma \cup N)^*$
4. a designated start symbol S

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CFG example

CFG's are also called phrase-structure grammars.
Equivalent to Backus-Naur Form (BNF).

- | | |
|---------------------------|-------------------------------------|
| 1. $S \rightarrow NP VP$ | 5. $NAME \rightarrow \text{Beavis}$ |
| 2. $VP \rightarrow V NP$ | 6. $V \rightarrow \text{ate}$ |
| 3. $NP \rightarrow NAME$ | 7. $ART \rightarrow \text{the}$ |
| 4. $NP \rightarrow ART N$ | 8. $N \rightarrow \text{cat}$ |
- CFG's are *powerful* enough to describe most of the structure in natural languages.
 - CFG's are *restricted* enough so that efficient parsers can be built.

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Derivations

- If the rule $A \rightarrow \beta \in P$, and α and γ are strings in the set $(\Sigma \cup N)^*$, then we say that $\alpha A \gamma$ **directly derives** $\alpha \beta \gamma$, or $\alpha A \gamma \Rightarrow \alpha \beta \gamma$
- Let $\alpha_1, \alpha_2, \dots, \alpha_m$ be strings in $(\Sigma \cup N)^*$, $m > 1$, such that

$$\alpha_1 \Rightarrow \alpha_2, \alpha_2 \Rightarrow \alpha_3, \dots, \alpha_{m-1} \Rightarrow \alpha_m,$$

then we say that α_1 **derives** α_m or $\alpha_1 \xRightarrow{*} \alpha_m$

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L_G

The language L_G generated by a grammar G is the set of strings composed of terminal symbols that can be derived from the designated start symbol S .

$$L_G = \{w | w \in \Sigma^*, S \xRightarrow{*} w\}$$

Parsing: the problem of mapping from a string of words to its parse tree according to a grammar G .

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General Parsing Strategies

Grammar	Top-Down	Bottom-Up
1. $S \rightarrow NP VP$	$S \rightarrow NP VP$	$\rightarrow \text{NAME ate the cat}$
2. $VP \rightarrow V NP$	$\rightarrow \text{NAME VP}$	$\rightarrow \text{NAME V the cat}$
3. $NP \rightarrow \text{NAME}$	$\rightarrow \text{Beav VP}$	$\rightarrow \text{NAME V ART cat}$
4. $NP \rightarrow \text{ART N}$	$\rightarrow \text{Beav V NP}$	$\rightarrow \text{NAME V ART N}$
5. $\text{NAME} \rightarrow \text{Beavis}$	$\rightarrow \text{Beav ate NP}$	$\rightarrow \text{NP V ART N}$
6. $V \rightarrow \text{ate}$	$\rightarrow \text{Beav ate ART N}$	$\rightarrow \text{NP V NP}$
7. $\text{ART} \rightarrow \text{the}$	$\rightarrow \text{Beav ate the N}$	$\rightarrow \text{NP VP}$
8. $N \rightarrow \text{cat}$	$\rightarrow \text{Beav ate the cat}$	$\rightarrow S$

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A Top-Down Parser

Input: CFG grammar, lexicon, sentence to parse

Output: yes/no

State of the parse: (*symbol list, position*)

1 The 2 old 3 man 4 cried 5

start state: ((S) 1)

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Grammar and Lexicon

Grammar:

1. $S \rightarrow NP VP$

2. $NP \rightarrow \text{art } n$

3. $NP \rightarrow \text{art adj } n$

4. $VP \rightarrow v$

5. $VP \rightarrow v NP$

Lexicon:

the: art

old: adj, n

man: n, v

cried: v

1 The 2 old 3 man 4 cried 5

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Algorithm for a Top-Down Parser

$PSL \leftarrow (((S) 1))$

1. *Check for failure.* If PSL is empty, return NO.
2. *Select the current state, C .* $C \leftarrow \text{pop}(PSL)$.
3. *Check for success.* If $C = (()) <\text{final-position}>$, YES.
4. *Otherwise, generate the next possible states.*
 - (a) $s_1 \leftarrow \text{first-symbol}(C)$
 - (b) If s_1 is a *lexical symbol* and next word can be in that class, create new state by removing s_1 , updating the word position, and adding it to PSL . (I'll add to front.)
 - (c) If s_1 is a *non-terminal*, generate a new state for each rule in the grammar that can rewrite s_1 . Add all to PSL . (Add to front.)

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Example

Current state	Backup states
1. $((S) 1)$	
2. $((NP VP) 1)$	
3. $((\text{art } n VP) 1)$	$((\text{art adj } n VP) 1)$
4. $((n VP) 2)$	$((\text{art adj } n VP) 1)$
5. $((VP) 3)$	$((\text{art adj } n VP) 1)$
6. $((v) 3)$	$((v NP) 3) ((\text{art adj } n VP) 1)$
7. $((()) 4)$	$((v NP) 3) ((\text{art adj } n VP) 1)$ Backtrack

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8. $((v NP) 3)$	$((\text{art adj } n VP) 1)$ leads to backtracking
...	
9. $((\text{art adj } n VP) 1)$	
10. $((\text{adj } n VP) 2)$	
11. $((n VP) 3)$	
12. $((VP) 4)$	
13. $((v) 4)$	$((v NP) 4)$
14. $((()) 5)$	$((v NP) 4)$
YES	

DONE!

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Problems with the Top-Down Parser

1. Only judges grammaticality.
2. Stops when it finds a single derivation.
3. No semantic knowledge employed.
4. No way to rank the derivations.
5. Problems with left-recursive rules.
6. Problems with ungrammatical sentences.

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Efficient Parsing

The top-down parser is terribly inefficient.

Have the first year Phd students in the computer science department take the Q-exam.

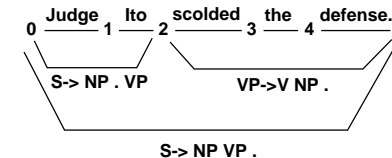
Have the first year Phd students in the computer science department taken the Q-exam?

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Chart Parsers

chart: data structure that stores partial results of the parsing process in such a way that they can be reused. The chart for an n -word sentence consists of:

- $n + 1$ **vertices**
- a number of **edges** that connect vertices



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Chart Parsing: The General Idea

The process of parsing an n -word sentence consists of forming a chart with $n + 1$ vertices and adding edges to the chart one at a time.

- Goal: To produce a complete edge that spans from vertex 0 to n and is of category S .
- There is no backtracking.
- Everything that is put in the chart stays there.
- Chart contains all information needed to create parse tree.

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Bottom-UP Chart Parsing Algorithm

Do until there is no input left:

1. If the agenda is empty, get next word from the input, look up word categories, add to agenda (as constituent spanning two positions).
2. Select a constituent from the agenda: constituent C from p_1 to p_2 .
3. Insert C into the chart from position p_1 to p_2 .
4. For each rule in the grammar of form $X \rightarrow C X_1 \dots X_n$, add an active edge of form $X \rightarrow C \circ X_1 \dots X_n$ from p_1 to p_2 .

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Bottom-up Chart Parser

Is it any less naive than the top-down parser?

1. Only judges grammaticality.[fixed]
2. Stops when it finds a single derivation.[fixed]
3. No semantic knowledge employed.
4. No way to rank the derivations.
5. Problems with ungrammatical sentences.[better]
6. Terribly inefficient.

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