Correction [4/25, 2pm]: typo in written problem 3: the label statement is "l : c", not "l : s".

Turn in all of the assignment using CMS. Turn in the written part in a file written.txt (or .doc, or .pdf), and submit your code in a file solve.ml. ..

## 1 Written part

- 1. Give a sample program fragment that produces different results for each of the following:
  - (a) Call-by-value and static scoping;
  - (b) Call-by-value and dynamic scoping;
  - (c) Call-by-name and dynamic scoping;
- 2. (a) Write a translation of exceptions into a language with if expressions, let constructs, pairs, and selection operators for the first and second components of the tuple;
  - (b) Extend the above translation to support constructs that pass values along with exceptions, e.g., " $e_1$  handle  $X(x) \Rightarrow e_2$ " and "throw X(e)".
- 3. We extend IMP with labels and goto statements:

The jump command goto l transfers the control at the point where the label command l:c occurs. For simplicity, assume that the label command occurs before the jump during the execution.

We'd like to model configurations as tuples (s, m, k, c) containing: a store s, a map m from labels to continuations, a current continuation k, and a currently executed command c. Continuations can be simply represented as commands, so m maps labels to commands, and k is some command c'. The evaluation of assignments is:

$$\begin{array}{ll} (s,m,c,x:=a) \to (s,m,c,x:=a') & \text{if } \langle a,s \rangle \to \langle a',s \rangle \\ (s,m,c,x:=v) \to (s[x\mapsto v],m,c,\texttt{skip}) \end{array}$$

Write the rules for skip,  $c_1$ ;  $c_2$ , if b then  $c_1$  else  $c_2$ , while b do c, l : c, and goto l. Note: make sure that while loops work correctly in the presence of goto's.

## 2 Programming part

In this part you will implement an interpreter for a mini-Prolog language called plog. A plog program consists of a set of rules (i.e., Horn clauses), followed by query that consists of a set of goals. A solution must satisfy all of the goals together. Here is an example program:

```
parent(bob, ann).
parent(bob, tom).
parent(bob, joe).
parent(amy, ann).
female(ann).
female(amy).
male(joe).
male(tom).
mother(X,Y) :- parent(X,Y), female(X).
father(X,Y) :- parent(X,Y), male(X).
? parent(X, Y), male(Y).
```

Terms and rules are represented using the following Ocaml types:

The program is a list of rules; and the query is a list of terms. Your job is to implement the function solve in file solve.ml:

solve: (rule list) -> (term list) -> (term -> term) list

For a list of rule  $\mathbf{r}$  and a set of query  $\mathbf{q}$ , the call **solve**  $\mathbf{r}$   $\mathbf{q}$  returns all the solutions to the query, i.e., the set of all unifiers that make  $\mathbf{q}$  feasible. If the query is not feasible, the returned list is empty. The main program in main.ml calls solve and displays all the solutions. For instance, the result for the above program is:

```
Found 2 solution(s).
Solution:
Y = joe
X = bob
Solution:
Y = tom
X = bob
```

For your convenience, we have provided a functional implementation of union-find structures (or unifiers); and a function match\_rule r crt rest that matches (via unification) a current goal crt with a rule r, for a remaining set of goals rest. The function match\_rule returns the new set of goals if unification succeeds; or raises FailUnify otherwise.

Each time a rule is matched, all of its variables must be alpha-renamed. The function fresh r renames all variables in rule r and returns the renamed rule. Different calls fresh r for the same rule r must yield different renamings.

To do: Implement the unification function unify and the renaming function fresh, and then use match\_rule to implement function solve. Turn in your solution in CMS, in file solve.ml.