Bali Types in your compiler

A how to

Things to consider

- Categories of Types:
 - Type "Signatures" (or Archetypes)
 - Primitive types
 - Integer
 - Boolean
 - Character
 - Function types
 - Class Types
 - (There are NOT in part 3)
 - Pointer types

A bali example:

```
class Entry

 int main()
        my int a1;
        my int* a2;
 char c(int a, boolean* b)
```

The symbol tables

- Not considering class symbol tables (part 3)
- Global Symbol Table:
 - Int type, Boolean type, Char type
 - ReadInt, ReadChar, ReadString
 - "Functions"
 - c
 - main

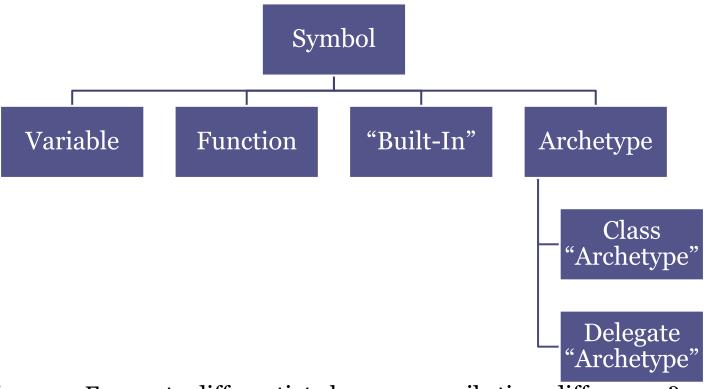
- Symbol Table for "main"
 - $a1 \rightarrow (int, SO = 2)$
 - $a2 \rightarrow (int^*, SO = 3)$
- Symbol Table for "c"
 - □ a -> (int, SO -2)
 - □ b -> (boolean*, SO= -1)

Symbol tables ARE mappings!

- What are we mapping from?
 - Strings
- What are we mapping to?
 - Symbols!

- What can we map?
 - Archetypes (eg: int, boolean, CustomClass)BUT NOT: pointers!!
 - Variables
 - Functions
 - "Built-In-Special" variables (eg: ReadChar)

Creating a type hierarchy



Suggestion: use Enums to differentiate known compile time differences? Also: Where do pointers fit in? Create a new class (eg: Type) that references an "Archetype" and has a field to denote the "pointer level".

What is a Delegate Archetype?

- Primitive types can be represented as Simple Archetypes (ie: by setting the type to some enum)
 - There is no other information that we need to know about primitive types
- For Functions/Delegates?
 - Return Type
 - Argument Types
 - Number of arguments

- For Classes
 - Field Members
 - Methods

Symbol Tables for functions

- Functions have their own symbol tables
- But the first block can not have variable names that conflict with function parameters
- Function and first block "share" a symbol table?

Solutions:

- All functions have a SINGLE initial block
- Have the function put the arguments into that initial block's symbol-table

Different "strokes" (symbol tables) for different "folks" (AST nodes)

- Different types of AST Nodes do different things
 - eg: They print out differentSaM code
- Different Symbol tables need to do different things
 - eg: For Blocks, the Symbol Table might "automatically" number the stack offset for variables.
 - You might need to use a different offset when doing classes
 - Or no numbering at all for the global symbol table?

Solutions:

- Can create different symbol tables that do different things when added to.
- Do you ever have to remove from a symbol table?
 - NO!!
- Also: Parameters are numbered differently in the symbol table than variables
- Only certain types of symbols can go in certain symbol tables (eg: "Built-Ins" can't go in a Block's Symbol Table)

Working with nested scopes

- AST Nodes that have scope (or create a dedicated symbol table)
 - Global (eg: Program)
 - Class (Not for this part)
 - Method (though might use the block's symbol table)
 - Block
- Since the Method uses the block's symbol table, we need only consider blocks
- Solutions:
 - Allocate variables "on demand"
 - If no return statement, just de-allocate variables at end of block (ADDSP)

- Solutions:
 - Can't conditionally allocate variables in a Block's scope, except via "Return Statement"
 - For Return (in nested block)
 - Keep track of variable allocations at the "method" level
 - Before jumping to "end of function", ADDSP the number of variables allocated at the method level
 - At the end of each block, subtract variable allocation of block from the method level
 - DON'T include parameters

AMS Demo and Questions