TREES

Lecture 12
CS2110 – Spring 2019

Announcements

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Submit P1 Conflict quiz on CMS by end of day Wednesday. We won't be sending confirmations; no news is good news. Extra time people will eventually get an email from Lacy. Please be patient.

Today's Topics in JavaHyperText

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- Search for "trees"
- Read PDFs for points 0 through 5: intro to trees, examples of trees, binary trees, binary search trees, balanced trees

Data Structures

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- **Data structure**
  - Organization or format for storing or managing data
  - Concrete realization of an abstract data type
- Operations
  - Always a tradeoff: some operations more efficient, some less, for any data structure
  - Choose efficient data structure for operations of concern

Example Data Structures

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Data Structure	add(val v)	get(int i)	contains(val v)
Array 2 1 3 0	$O(n)$	$O(1)$	$O(n)$
Linked List 2→1→3→0	$O(1)$	$O(n)$	$O(n)$

```
add(v): append v
get(i): return element at position i
contains(v): return true if contains v
```

Tree

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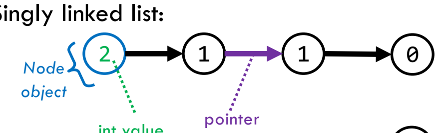
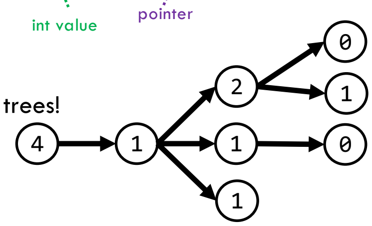
Singly linked list:



Today: trees!



Trees

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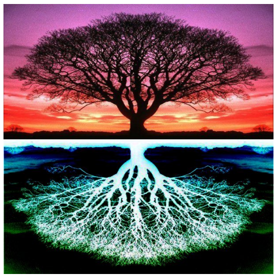

In CS, we draw trees “upside down”

Tree Overview

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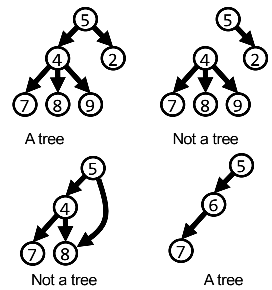
Tree: data structure with nodes, similar to linked list

- Each node may have zero or more *successors* (children)
- Each node has exactly one *predecessor* (parent) except the *root*, which has none
- All nodes are reachable from *root*

A tree or not a tree?



Tree Terminology (1)

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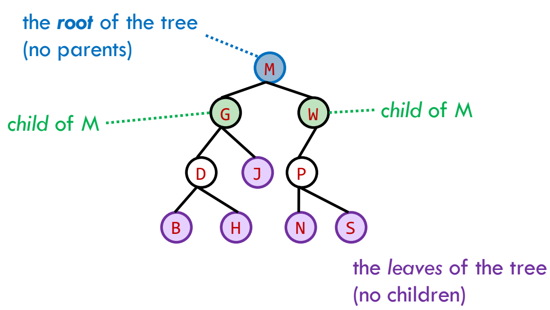

Tree Terminology (2)

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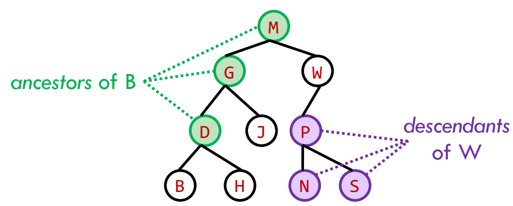

Tree Terminology (3)

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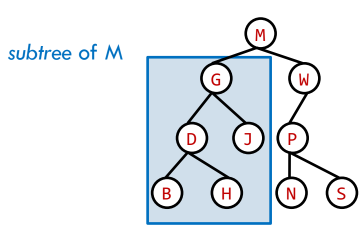

Tree Terminology (4)

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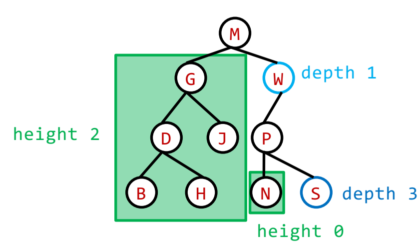
A node’s ***depth*** is the length of the path to the root.

A tree’s (or subtree’s) ***height*** is the length of the longest path from the root to a leaf.



Tree Terminology (5)

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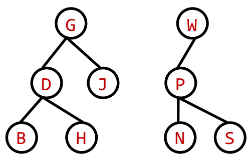
Multiple trees: a ***forest***



General vs. Binary Trees

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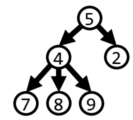
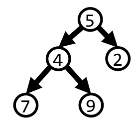
General tree: every node can have an arbitrary number of children



General tree

Binary tree: at most two children, called *left* and *right*



Binary tree

…often "tree" means binary tree

Demo

Binary trees were in A1!

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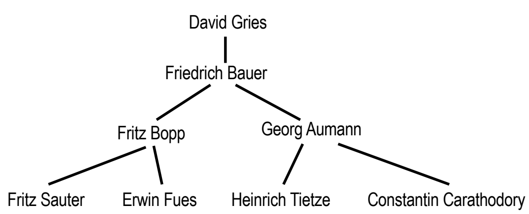
You have seen a binary tree in A1.
A PhD object has one or two advisors.
(Note: the advisors are the "children".)



Special kinds of binary trees

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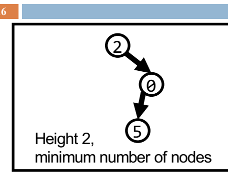
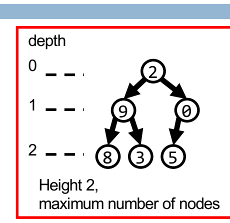


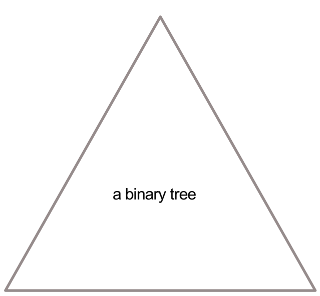
Max # of nodes at depth d: 2^d

If height of tree is h:
min # of nodes: $h + 1$
max #of nodes: (Perfect tree)
$2^0 + \ldots + 2^h = 2^{h+1} - 1$

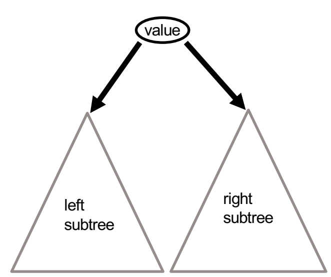


Complete binary tree
Every level, except last, is completely filled, nodes on bottom level as far left as possible. No holes.

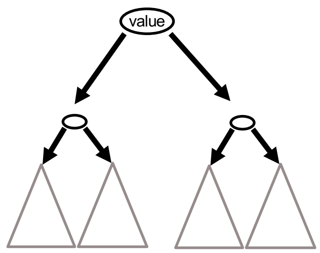
Trees are recursive



Trees are recursive

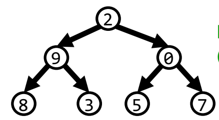


Trees are recursive



Trees are recursive

Left subtree, which is also a binary tree



Right subtree (also a binary tree)

Trees are recursive

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A binary tree is either `null`

or an object consisting of a value, a left binary tree, and a right binary tree.

A Recipe for Recursive Functions

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Base case:
    If the input is "easy," just solve the problem directly.

Recursive case:
    Get a smaller part of the input (or several parts).
    Call the function on the smaller value(s).
    Use the recursive result to build a solution for the full input.

A Recipe for Recursive Functions on Binary Trees

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Base case: an empty tree (null), or possibly a leaf
    If the input is "~~easy~~," just solve the problem directly.

Recursive case:
    ~~Get a smaller part of the input (or several parts).~~
    Call the function on ~~the smaller value(s)~~ each subtree
    Use the recursive result to build a solution for the full input.

Demo

Comparing Searches

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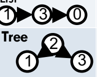
Data Structure	add(val v)	get(int i)	contains(val v)
Array 2 1 3 0	$O(n)$	$O(1)$	$O(n)$
Linked List 2→1→3→0	$O(1)$	$O(n)$	$O(n)$
Binary Tree 			$O(n)$

Node could be *anywhere* in tree

Binary search on arrays: O(log n)
Requires invariant: array sorted
...analogue for trees?

Binary Search Tree (BST)

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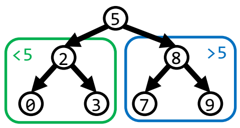
A *binary search tree* is a binary tree with a **class invariant**:

- All nodes in the left subtree have values that are less than the value in that node, and
- All values in the right subtree are greater.

(assume no duplicates)



Demo

Binary Search Tree (BST)

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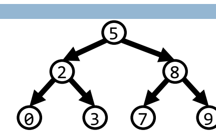

Contains:

- Binary tree: two recursive calls: O(n)
- BST: one recursive call: O(height)

BST Insert

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To insert a value:

- Search for value
- If not found, put in tree where search ends

Example: Insert month names in chronological order as Strings, (Jan, Feb…). BST orders Strings alphabetically (Feb comes *before* Jan, etc.)

BST Insert

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insert: January

BST Insert

29

insert: February

BST Insert

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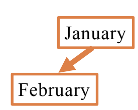
insert: March



BST Insert

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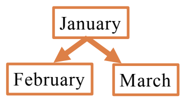
insert: April…



BST Insert

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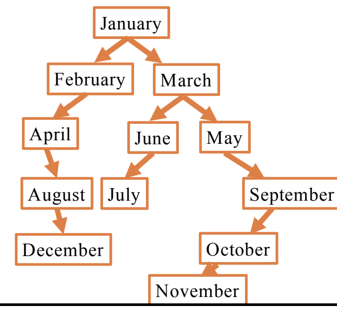

Comparing Data Structures

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Data Structure	add(val x)	get(int i)	contains(val x)
Array 2 1 3 0	$O(n)$	$O(1)$	$O(n)$
Linked List 2 1 3 0	$O(1)$	$O(n)$	$O(n)$
Binary Tree 1 2 3			$O(n)$
BST 2 1 3	$O(height)$		$O(height)$

How big could height be?

Worst case height

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Insert in alphabetical order…

Worst case height

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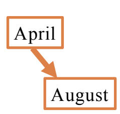
Insert in alphabetical order…



Worst case height

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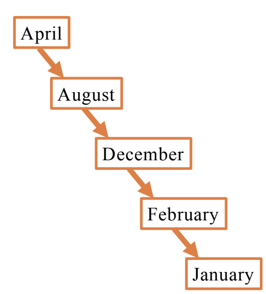
Insert in alphabetical order…



Tree degenerates to list!

Need Balance

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- Takeaway: BST search is O(n) time
  - Recall, big O notation is for worst case running time
  - Worst case for BST is data inserted in sorted order
- **Balanced binary tree:** subtrees of any node are about the same height
  - In balanced BST, search is O(log n)
  - Deletion: tricky! Have to maintain balance
  - [Optional] See JavaHyperText "Extensions to BSTs"
  - Also see CS 3110