



1

ASTS, GRAMMARS, PARSING, TREE TRAVERSALS

Lecture 14
CS2110 – Fall 2018

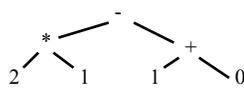
Announcements

- Today: The last day to request prelim regrades
- Assignment A4 due next Thursday night. Please work on it early and steadily. Watch the two videos on recursion on trees before working on A4!
- Next week's recitation. Learn about interfaces Iterator and Iterable. There will be 15 minutes of videos to watch. Then, in recitation, you will fix your A3 so that a foreach loop can be used on it.

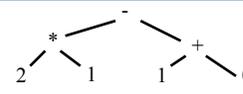
```
DLL<Integer> d= new DLL<Integer>();
...
for (Integer i : d) { ... }
```

Expression Trees

we can draw a **syntax tree** for the Java expression `2 * 1 - (1 + 0)`.



Pre-order, Post-order, and In-order

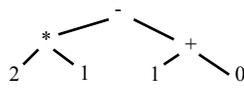


Pre-order traversal:

1. Visit the root
2. Visit the left subtree (in pre-order)
3. Visit the right subtree

- * 2 1 + 1 0

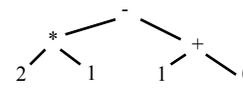
Pre-order, Post-order, and In-order



Pre-order traversal **- * 2 1 + 1 0**
 Post-order traversal **2 1 * 1 0 + -**

1. Visit the left subtree (in post-order)
2. Visit the right subtree
3. Visit the root

Pre-order, Post-order, and In-order



Pre-order traversal **- * 2 1 + 1 0**
 Post-order traversal **2 1 * 1 0 + -**
 In-order traversal **2 * 1 - 1 + 0**

1. Visit the left subtree (in-order)
2. Visit the root
3. Visit the right subtree

Pre-order, Post-order, and In-order

Pre-order traversal **- * 2 1 + 1 0**

Post-order traversal **2 1 * 1 0 + -**

In-order traversal **(2 * 1) - (1 + 0)**

To avoid ambiguity, add parentheses around subtrees that contain operators.

In Defense of Postfix Notation

- Execute expressions in postfix notation by reading from left to right.
- Numbers: push onto the stack.
- Operators: pop the operands off the stack, do the operation, and push the result onto the stack.

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13

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14

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15

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In about 1974, Gries paid \$300 for an HP calculator, which had some memory and used postfix notation! Still works.

a.k.a. "reverse Polish notation"

In Defense of Prefix Notation

16

- Function calls in most programming languages use prefix notation: like add(37, 5).
- Some languages (Lisp, Scheme, Racket) use prefix notation for *everything* to make the syntax simpler.

```
(define (fib n)
  (if (<= n 2)
      1
      (+ (fib (- n 1)) (fib (- n 2)))))
```

Determine tree from preorder and postorder

18

Suppose inorder is **B C A E D**
preorder is **A B C D E**

Can we determine the tree uniquely?

Determine tree from preorder and postorder

19

Suppose inorder is **B C A E D**
preorder is **A B C D E**

Can we determine the tree uniquely?

What is the root? preorder tells us: **A**

What comes before/after root **A**? Inorder tells us:
Before : **B C**
After: **E D**

Determine tree from preorder and postorder

20

Suppose inorder is **B C A E D**
preorder is **A B C D E**

The root is **A**.

Left subtree contains **B C** Right subtree contains **E D**

Now figure out left, right subtrees *using the same method*.

From the above:

For left subtree

inorder is: **B C**

preorder is: **B C**

root is: **B**

Right subtree: **C**

For right subtree:

inorder is: **E D**

preorder is: **D E**

root is: **D**

left subtree: **E**

Expression trees: in code

21

```
public interface Expr {
    String inorder(); // returns an inorder representation
    int eval(); // returns the value of the expression
}
```

```
public class Int implements Expr {
    private int v;
    public int eval() { return v; }
    public String inorder() {
        return "" + v + "";
    }
}

public class Sum implements Expr {
    private Expr left, right;
    public int eval() {
        return left.eval() + right.eval();
    }
    public String inorder() {
        return "(" + left.infix() +
            "+" + right.infix() + ")";
    }
}
```

Grammars

22

- The cat ate the rat.
- The cat ate the rat slowly.
- The small cat ate the big rat slowly.
- The small cat ate the big rat on the mat slowly.
- The small cat that sat in the hat ate the big rat on the mat slowly, then got sick.

- Not all sequences of words are sentences:
The ate cat rat the
- How many legal sentences are there?
- How many legal Java programs are there?
- How can we check whether a string is a Java program?

Grammars

23

A **grammar** is a set of rules for generating the valid strings of a language.

Read → as “may be composed of”

- Sentence → Noun Verb Noun
- Noun → goats
- Noun → astrophysics
- Noun → bunnies
- Verb → like
- Verb → see

Grammars

24

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 - Verb → see
- Sentence

Grammars

25

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- Noun Verb Noun

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26

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28

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Grammars

30

A **grammar** is a set of rules for generating the valid strings of a language.

Sentence → Noun Verb Noun
 Noun → goats
 Noun → astrophysics
 Noun → bunnies bunnies like astrophysics
 Verb → like goats see bunnies
 Verb → see ... (18 sentences total)

- The words **goats**, **astrophysics**, **bunnies**, **like**, **see** are called *tokens* or *terminals*
- The words **Sentence**, **Noun**, **Verb** are called *nonterminals*

A recursive grammar

31

Sentence → Sentence and Sentence
 Sentence → Sentence or Sentence
 Sentence → Noun Verb Noun
 Noun → goats
 Noun → astrophysics
 Noun → bunnies
 Verb → like bunnies like astrophysics
 → see goats see bunnies
 bunnies like goats and goats see bunnies
 ... (infinite possibilities!)

The recursive definition of **Sentence** makes this grammar infinite.

Grammars for programming languages

34

A grammar describes every possible legal program.

You could use the grammar for Java to list every possible Java program. (It would take forever.)

A grammar also describes how to “parse” legal programs.

The Java compiler uses a grammar to translate your text file into a syntax tree—and to decide whether a program is legal.

docs.oracle.com/javase/8/docs/spotlights/html/jls-2.html#jls-2.3

docs.oracle.com/javase/8/docs/spotlights/html/jls-19.html

Grammar for simple expressions (not the best)

35

$E \rightarrow \text{integer}$

$E \rightarrow (E + E)$

Simple expressions:

- An E can be an integer.
- An E can be '(' followed by an E followed by '+' followed by an E followed by ')'

Set of expressions defined by this grammar is a recursively-defined set

- Is language finite or infinite?
- Do recursive grammars always yield infinite languages?

Some legal expressions:

- 2
- (3 + 34)
- ((4+23) + 89)

Some illegal expressions:

- (3
- 3 + 4

Tokens of this grammar:

(+) and any integer