



Lecture 26: More on Algorithms for Sorting

CS 1110

Introduction to Computing Using Python

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Announcements

- Discussion sections this week
 - First 10 minutes dedicated to getting your started on A6
 - Remaining time is office hour for your A6/Prelim 2 questions
- Final Exam on May 21st 1:30-4pm. Your assigned exam session (in-person or online) is shown in CMS. Submit a “regrade request” in CMS by May 12 if you have a *legitimate* reason for requesting a change. If you have an exceptional circumstance for switching from in-person to online, you must upload to CMS your college’s approval of your modality change.

More Announcements

- A6 due on Friday
 - Remember academic integrity
- Expected release dates of solutions and feedback
 - A5 solutions: Wed May 12
 - A4 grades and feedback: Thurs May 13
 - A6 solutions: Tues May 18
 - A5 grades and feedback: Thurs May 20
 - Final exam grades and feedback: Tues May 25
 - A6 grades and feedback: Fri May 28

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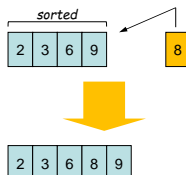
Algorithms for Sorting

- Well known algorithms
 - focus on reviewing programming constructs (**while** loop) and analysis
 - will not use built-in methods such as **sort**, **index**, **insert**, etc.
- Today we’ll discuss **merge sort** and compare it to **insertion sort**, which we discussed last lecture
- More on the topic in next course, CS 2110!

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The Insertion Process of Insertion Sort

- Given a sorted list x , insert a number y such that the result is sorted
- Sorted: arranged in ascending (small to big) order



We’ll call this process a “push down,” as in push a value down until it is in its sorted position

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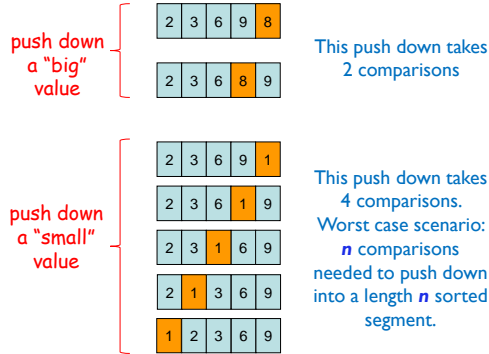
Algorithm Complexity

- Count the number of comparisons needed
- In the worst case, need i comparisons to push down an element in a sorted segment with i elements.

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How much work is a push down?



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Algorithm Complexity (Q)

```
def swap(b, h, k):
    :
def push_down(b, k):
    while k > 0 and b[k-1] > b[k]:
        swap(b, k-1, k)
        k = k-1
def insertion_sort(b):
    for i in range(1, len(b)):
        push_down(b, i)
```

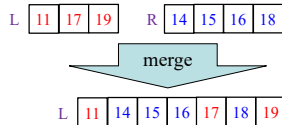
Count (approximately) the number of **comparisons** needed to sort a list of length n

- A. ~ 1 comparison
- B. ~ n comparisons
- C. ~ n^2 comparisons
- D. ~ n^3 comparisons
- E. I don't know

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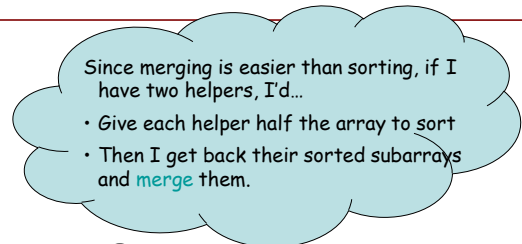
Which algorithm does Python's sort use?

- Recursive algorithm that runs much faster than insertion sort for the same size list (when the size is big!)
- A variant of an algorithm called "merge sort"
- Based on the idea that sorting is hard, but "merging" two *already sorted* lists is easy.



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Merge sort: Motivation



What if those two helpers each had two sub-helpers?

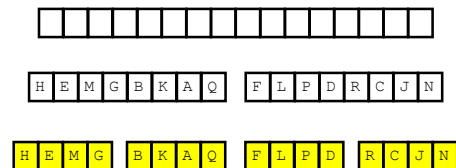
And the sub-helpers each had two sub-sub-helpers? And...

Subdivide the sorting task



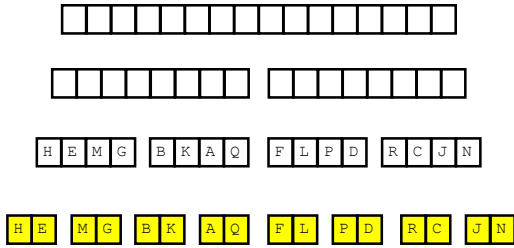
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Subdivide again



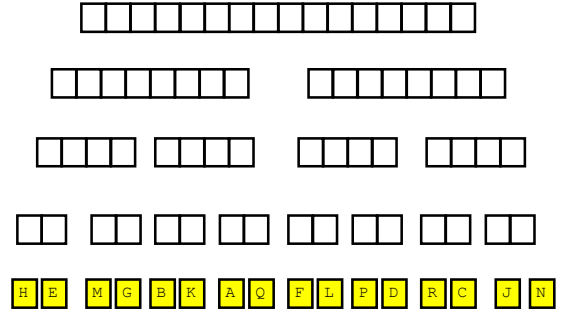
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And again

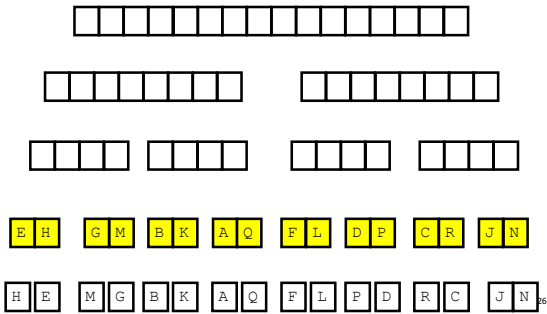


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And one last time

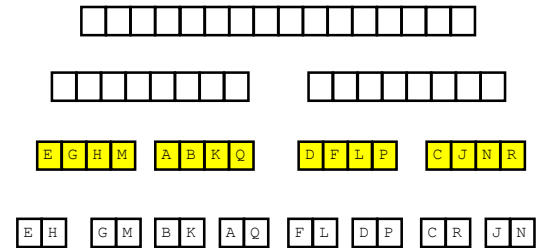


Now merge



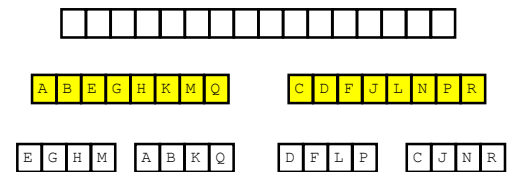
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And merge again



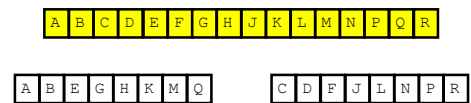
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And again



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And one last time



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Done!



```
def mergeSort(li):
    """Sort list li using Merge Sort"""
    if len(li) > 1:
        # Divide into two parts
        mid= len(li)//2
        left= li[:mid]
        right= li[mid:]

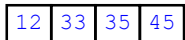
        # Recursive calls
        mergeSort(left)
        mergeSort(right)

        # Merge left & right back to li
    ...
```

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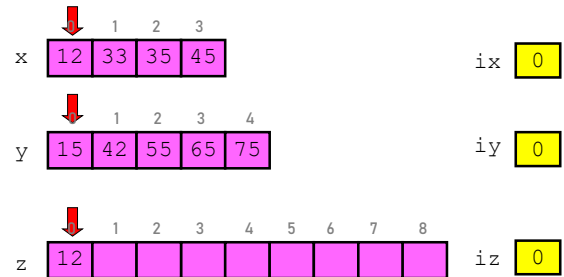
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The central sub-problem is the merging of two sorted lists into one single sorted list



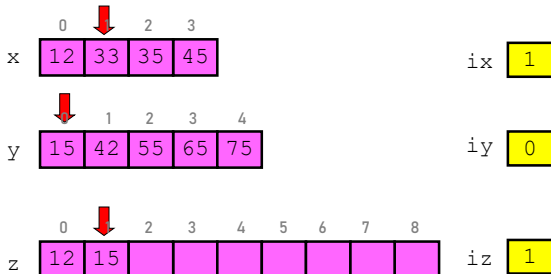
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Merge



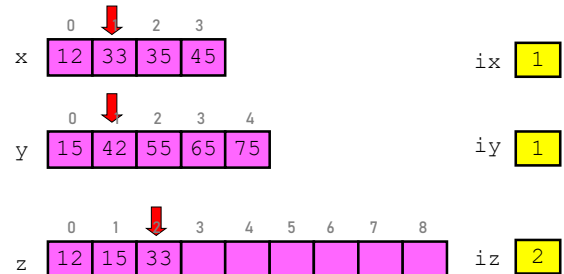
$ix < 4$ and $iy < 5 \rightarrow x(ix) \leq y(iy)$ YES

Merge



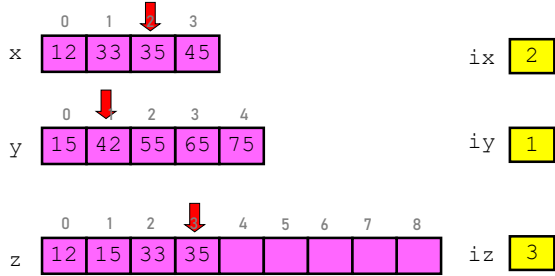
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Merge



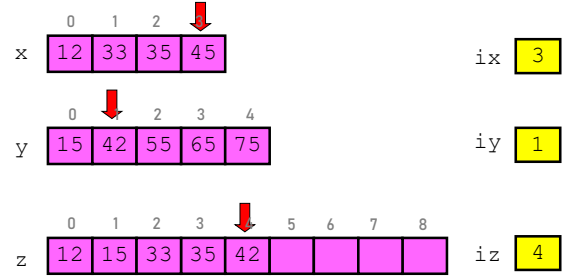
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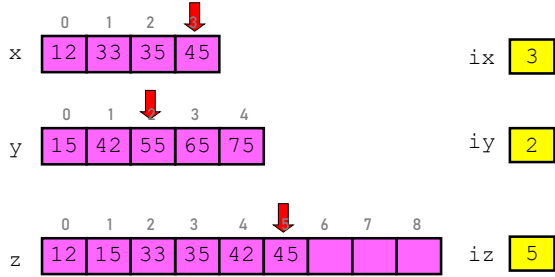
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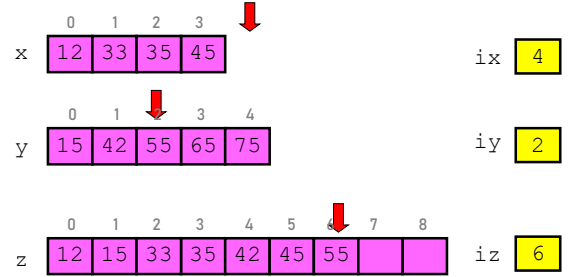
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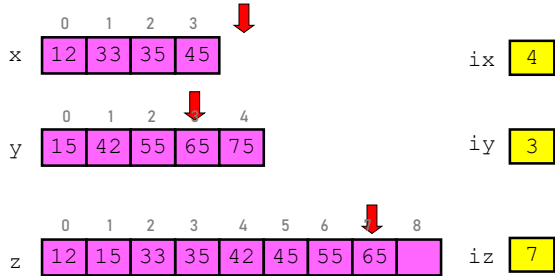
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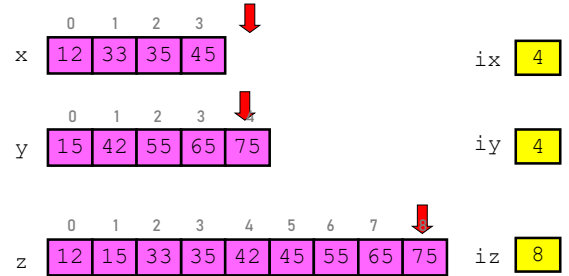
ix at 4 \rightarrow take $y(iy)$

Merge

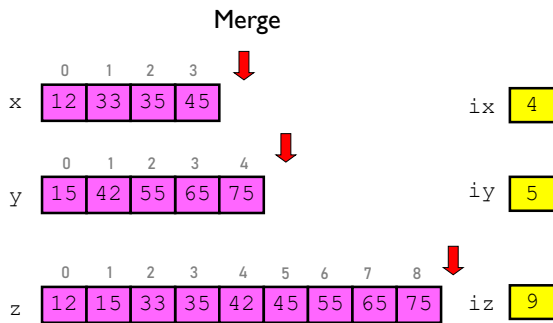


$iy < 5 \rightarrow$ take $y(iy)$

Merge



$iy < 5 \rightarrow$ take $y(iy)$



```
# Given lists x and y and list z, which has
# the combined length of x and y...
nx = len(x); ny = len(y)
```

```
ix = 0; iy = 0; iz = 0;
while ix<nx and iy<ny:
    if x[ix] <= y[iy]:
        z[iz]= x[ix]; ix=ix+1
    else:
        z[iz]= y[iy]; iy=iy+1
    iz=iz+1
```

```
while ix<nx # copy any remaining x-values
    z[iz]= x[ix]; ix=ix+1; iz=iz+1
```

```
while iy<ny # copy any remaining y-values
    z[iz]= y[iy]; iy=iy+1; iz=iz+1
```

How do merge sort and insertion sort compare?

- Insertion sort: (worst case) makes *i* comparisons to insert an element in a sorted array of *i* elements. For an array of length *n*:

_____ for big *n*

- Merge sort: _____

```
def mergeSort(li):
    """Sort list li using Merge Sort"""
    if len(li) > 1:
        # Divide into two parts
        mid= len(li)/2
        left= li[:mid]
        right= li[mid:]

        # Recursive calls
        mergeSort(left)
        mergeSort(right)

        # Merge left & right back to li
        ...
```

All the comparisons between list elements are done during merge

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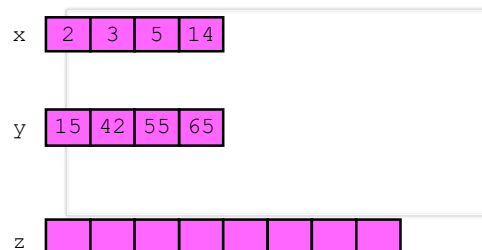
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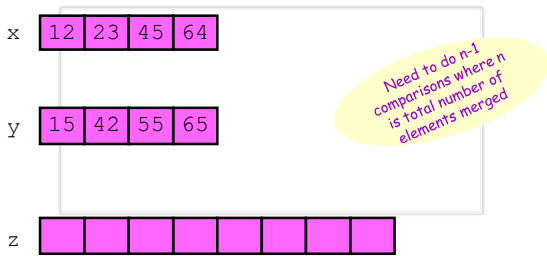
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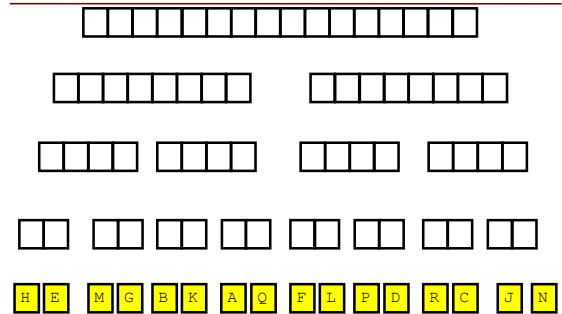
Merge – best case scenario



Merge – worst case scenario



Merge sort: about $\log_2(n)$ “levels”;
about n comparisons each level



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How do merge sort and insertion sort compare?

- Insertion sort: (worst case) makes i comparisons to insert an element in a sorted array of i elements. For an array of length n :
 $1+2+\dots+(n-1) = n(n-1)/2$, say n^2 for big n
- Merge sort: $n \cdot \log_2(n)$ comparisons
- Should we always use merge sort then? Python actually uses a variant that combines merge sort and insertion sort!

Order of magnitude difference