

Lecture 15: Loop invariants

- How to develop loops, guaranteeing correctness
- Also: finishing heaps, start sorting

Announcements

- P3 extension (again)
- No T/W discussion
- Two surveys (sorry!)

Heap implements Priority Queue:

- insert (value, priority)
- removeMax()
- findMax()

Heap implementation:

binary tree satisfying:

- ① each node is larger than its children
- ② the tree is full.

} class invariants

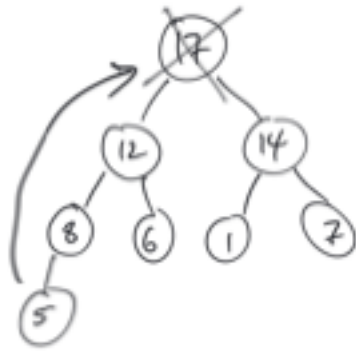
• max is at root

• to insert: put new val in last pos
- swap with parent & repeat as necessary to satisfy inv. ①.

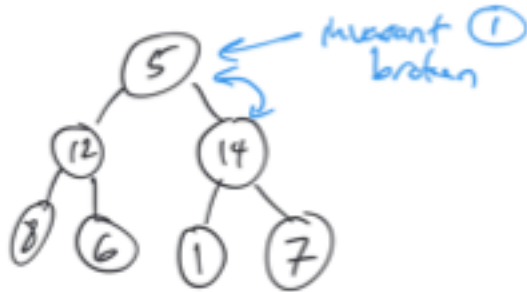


• removeMax():

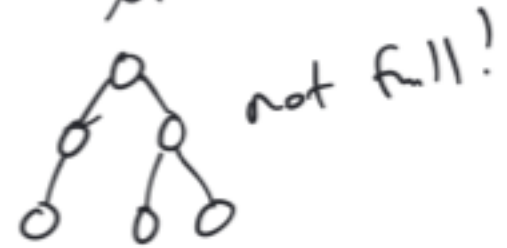
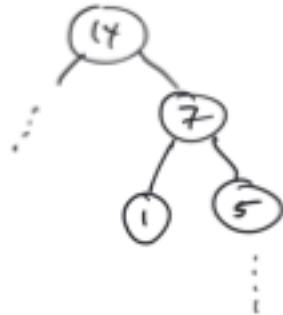
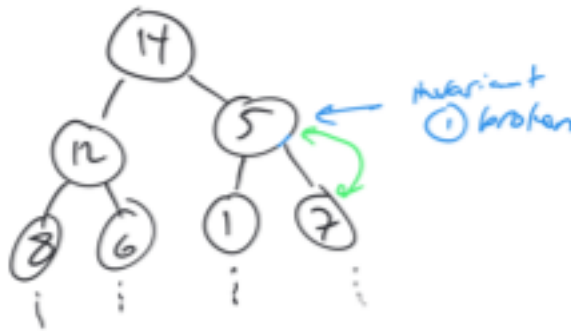
- ① replace the root with the last leaf



- ② swap root with larger of its children



repeat until invariant holds



not full!

loop invariant: something you know is true in each iteration of a loop.

Example: remove from heap.

- save root value to return
- swap last value in for root (remove it).

To develop a loop (4 loopy questions):

- INIT: does it start right?

set variables so that precondition makes invariant true

- TERMINATE: does it end right?

create loop guard that, with invariants, guarantees postcondition

- PROGRESS: will it end?

start loop body by making progress towards term. cond.

- PRESERVATION: is it invariant?

finish loop body by ensuring that invariant still holds after