

## CS100J Fall 2003 Answers to Exercises on loops

- E1(a)** `k= 2; x= 2; // (because of the conjunct  $2 \leq k$  in the invariant, we can't set k to 1)  
// invariant: P1:  $2 \leq k \leq 10$  and x is the product of 2..k  
while (k != 10) {  
    x= x * (k+1)  
    k= k+1;  
}  
// postcondition R is: x is the product of 2..10`
- E1(b)** `k= 2; x= 1; // (the product of no values is 1; the sum of no values is 0)  
// invariant P2:  $2 \leq k \leq 11$  and x is the product of 2..(k - 1)  
while (k != 11) {  
    x= x*k;  
    k= k + 1;  
}  
// postcondition R is: x is the product of 2..10`
- E1(c)** `k= 10; x= 10;  
// invariant: P3:  $2 \leq k \leq 10$  and x is the product of k..10  
while (k != 2) {  
    x= x * (k-1);  
    k= k-1;  
}  
// postcondition R is: x is the product of 2..10`
- E1(d)** `k= 10; x= 1;  
// invariant: P4:  $1 \leq k \leq 10$  and x is the product of (k + 1)..10  
while (k != 1) {  
    x= x * k;  
    k= k-1;  
}  
// postcondition R is: x is the product of 2..10`
- E2(a)** Note: the conjunct !b in the loop condition is not necessary. It was added later, as an afterthought, for the following reason. The repetend never falsifies b, so once b becomes true, meaning that some integer in the range divides n, the loop can terminate. This holds for all four subexercises.  
`k= first-1; b= false;  
// P1:  $first - 1 \leq k \leq last$  and b = "n is divisible by an integer in first..k"  
while (!b && k != last) {  
    if ( n % k == 0)  
        b= true;  
    k= k+1;  
}  
// postcondition R: b = "n is divisible by an integer in first..last"`
- E2(b)** `k= first; b= false;  
// invariant: P2: b = "n is divisible by an integer in first..(k - 1)"  
while (!b && k-1 != last) {  
    if (n % k == 0)  
        b= true;  
    k= k+1;  
}  
// postcondition R: b = "n is divisible by an integer in first..last"`
- E2(c)** `k= last + 1;  
b= false;  
// invariant: P3: b = "n is divisible by an integer in k..last"`

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while (!b && k != first) {
    if (n%(k-1) == 0)
        b= true;
    k= k - 1;
}
// postcondition R: b = "n is divisible by an integer in first..last"

E2(d) k= last;
b= false;
// invariant: P4: b = "n is divisible by an integer in k+1..last"
while (!b && k+1 != first) {
    if (n % k == 0)
        b= true;
    k= k - 1;
}
// postcondition R: b = "n is divisible by an integer in first..last"

E3. //Precondition: n > 0
int k = 0; int b = 1;
// invariant:  $1 \leq 2^{**}k \leq n$  and  $b = 2^{**}k$ 
while (2*b <= n){
    b= 2*b;
    k= k + 1;
}
// postcondition:  $1 \leq 2^{**}k \leq n < 2^{**}(k+1)$ 

E4. // precondition:  $x \geq 0$  and  $y > 0$ 
int q= 0;
int r= x;
// invariant:  $x = y * q + r$ 
while (r > y){
    q= q + 1;
    r= r - y;
}

E5. //precondition:  $x > 0$  and  $y > 0$  are integers
int b= x;
int c= y;
// invariant:  $b \text{ gcd } c = x \text{ gcd } y$ 
while ( b != c){
    /* use whichever property of gcd (given in the exercise description
       that makes progress (decreases b or c) while keeping b and c positive
       and maintain. */
    if (b > c)b= b - c;
    else    c= c - b;
}

E6. // precondition: t is a String and not null.
StringBuffer s= new StringBuffer(t);
int k = 0;
//invariant: s[0..(k-1)] contains no vowels and  $k \neq s.length() + 1$ 
while (k != s.length()){
    // c= the character at index k, converted to lower case;
    char c= s.charAt(k);
    c= Character.toLowerCase(c);

    // Make progress here by either decreasing s.length()
    // (by removing a character) or increasing k.

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        if (c == 'a' || c == 'e' || c == 'i' || c == 'o' || c == 'u'){
            s.deleteCharAt(k);
        } else {
            k= k + 1;
        }
    }
    t= s.toString();
E7. //precondition: s1.length() = s2.length()
    int k= 0;
    boolean areComplements= true;
    // invariant: areComplements = "s1[0..(k-1)] is the DNA complement of s2[0..(k-1)]"
    while (areComplements && k != s1.length()){
        char c1= s1.charAt(k);
        char c2= s2.charAt(k);
        // code below results in areComplements = 'c1 and c2 are DNA complements'
        if (c1 == 'A' & (c2 != 'T')) { areComplements= false; }
        if (c1 == 'C' & (c2 != 'G')) { areComplements= false; }
        if (c1 == 'G' & (c2 != 'C')) { areComplements= false; }
        if (c1 == 'T' & (c2 != 'A')) { areComplements= false; }
        k= k+1;
    }
}
E8. String s_comp= "";
    // invariant: s_comp is the DNA complement of s[0..k-1]
    for (int k= 0; k < s.length(); k= k + 1 ) {
        if (s.charAt(k) == 'A') { s_comp= s_comp + 'T'; }
        else if (s.charAt(k) == 'T') { s_comp= s_comp + 'A'; }
        else if (s.charAt(k) == 'G') { s_comp= s_comp + 'C'; }
        else { s_comp= s_comp + 'G'; }
    }
    // s_comp is the DNA complement of a s
E9. // precondition: n > 0
    int b = 0;
    int a = 1;
    int i = 1;
    // invariant: a = f[i] and b = f[i-1]
    while (i < n) {
        temp= a + b;
        b= a;
        a= temp;
        i= i + 1;
    }
    // postcondition: i = n (and, therefore, a = fn)
E10. sum = 0;
    // invariant: sum = sum of integers that are already read from file
    while (in.available) {
        i= in.readInt();
        sum= sum + i;
    }
    // postcondition: sum = sum of the integers in the file
E11. int countOdd= 0;
    int countEven= 0;

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// in: countOdd and countEven contain the number of odd and even integers already read from file
while (in.available) {
    int i= in.readInt();
    if (i % 2 == 0)
        countEven= countEven + 1;
    else
        countOdd= countOdd + 1;
}

// postcondition: countOdd and countEven contain the number of odd and even integers in the file

E12. // precondition: n >= 0;
int i = 0;
// invariant: balance contains the balance in the account after i years
while (i < n) {
    balance= balance + balance * rate;
    i= i + 1;
}

// postcondition: balance contains the balance in the account after n years

E13. double e= 1;
int k= 0;
double tk= 1;
// invariant: e = 1/0! + 1/1! + 1/2! + 1/3! + 1/4! + ... + 1/k! and tk = 1/k!
while (tk >= 1E-14) {
    k= k+1;
    tk= tk/k;
    e= e + tk;
}

E14. The loop below required 20001 iterations to find the approximation 3.141597653564762 to pi =
3.141592653589793. That's far too long, and this is not a good way to calculate pi.

int k= 0;
double t = 4;
double pi= t;
int sgn= 1;
// invariant: pi = 4/1 - 4/3 + 4/5 - 4/7 + 4/9 - ... + (-1)**k*4/(2k+1) and
//          t = 4/(2*k+1) and
//          sgn = (-1)**k
while (t >= .00001) {
    k= k+1;
    t= 4.0/(2*k+1);
    sgn= -sgn;
    pi= pi + sgn*t;
}

E15. This loop is preferable to that of E14 because it took only 10 iterations to stop with the same stopping
conditions

double c= 2.0*Math.sqrt(3);
int k= 0;
double term= c;
int t= 1;
double pi= c;
// inv: pi = c/(1*3**0) - c/(3*3**1) + c/(5*3**2) - c/(7*3**3) + ...
//          + (-1)**k * c * / ((2*k+1) * 3**k) and
//          t = 3**k and
//          term = c * / ((2*k+1) * 3**k)

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while (term >= .00001) {
    k= k+1;
    t= t*3;
    term= c/((2*k+1)*t);
if (k%2 == 1) pi= pi - term;
else pi= pi + term;
}

E16. int ndarts= 10000; // number of darts to throw
int k= 0;
int nhits= 0;
java.util.Random rand = new java.util.Random(System.currentTimeMillis());
/* invariant: nhits = number of hits after k darts thrown */
while (k < ndarts) {
double x = 2 * rand.nextDouble() - 1;
double y = 2 * rand.nextDouble() - 1;
if (x*x + y*y <= 1)
    nhits= nhits + 1;
    k= k + 1;
}
/* postcondition: nhits = no. of hits after darts thrown and k = ndarts */
double pi = 4.0 * nhits / ndarts;

E17. int n= 0;
// inv: n is the number of times 'a' occurs in s[0..k-1]
for (int k= 0; k != s.length(); k= k+1) {
    if (s.charAt(k) == 'a') {
        n= n+1;
    }
}
// postcondition: n is the number of times 'a' occurs in s

E18. This solution looks only for lowercase vowels
int numVowels= 0;
int i= -1;
// invariant: numVowels is the number of vowels in s[0..i]
while (i != s.length()-1) {
    i= i + 1;
    if (s.charAt(i) == 'a' || s.charAt(i) == 'e' || s.charAt(i) == 'i' ||
        s.charAt(i) == 'o' || s.charAt(i) == 'u')
        numVowels= numVowels + 1;
}
// postcondition: numVowels is the number of vowels in s[0..s.length()-1]

E19. int adjEqChars= 0;
int i= 0;
// invariant: adjEqChars is the number of adjacent equal characters in s[0..i]
while (i < s.length()-1) {
    if (s.charAt(i) == s.charAt(i+1))
        adjEqChars= adjEqChars + 1;
    i= i + 1;
}
// postcondition: adjEqChars is the number of adjacent equal characters in s[0..s.length()-1]

E20. int i= 0;
// invariant: the characters in s[0..i-1] are in descending order
while (i < s.length()-1 && s.charAt(i) >= s.charAt(i+1)) {
    i= i + 1;
}

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}
// postcond.: the chars in s[0..i-1] are in descending order but the next char, if it exists, is not
E21. int i= s.length()-1;
// invariant: s[i+1..s.length()-1] is all blanks
while (i >= 0 && s.charAt(i) == ' ') {
    i = i - 1;
}
// postcondition: i is the number of blanks at the end of s

E22. int h= 0;
int k= s.length() - 1;
// invariant: s[0..h-1] is the reverse of s[k+1..s.length()-1]
while (h < k && s.charAt(h) == s.charAt(k)) {
    h = h+1; k= k-1;
}
// postcondition: s[0..h-1] is reverse of s[k+1..s.length()-1] and
// either h >= k or s.charAt(h) == s.charAt(k)
boolean b= (h >= k); // b = "s is a palindrome"

E23. int k = 0;
boolean b= true;
// inv: b = "every character in s[0..k-1] has the same char next to it in s"
while (b && k < s.length()) {
// Set before to "s[k] has the same character before it"
boolean before= k != 0 && s.charAt(k) == s.charAt(k-1);

// set after to "s[k] has the same character after it"
boolean after= k+1 < s.length() && s.charAt(k) == s.charAt(k+1);

    b= before || after;
    k = k + 1;
}

E24. int k=1;
boolean b = true;
// invariant: b = "s[c] of s[0..k-1] is a digit for any c that is a power of two"
while (b & k < s.length()) {
    b = b && Character.isDigit(s.charAt(k));
    k = 2 * k;
}
// postcondition: b = s[c] is a digit for any c that is a power of two

E25. boolean b= s.length() == t.length();
// invariant: b = "s and t have the same length and s[0..k-1] = t[0..k-1]"
for (int k= 0; b && k < s.length(); k= k+1) {
    if (s.charAt(k) != t.charAt(k))
        b= false;
}
// postcondition: b = (s[]= t)

E26. String t= "";
int k= 0;
// invariant: t is s[0..k-1] but with twins added and
// if s[k-1] has a twin, it is s[k-2]
while (k != s.length()) {
    // Append two copies of s[k] to t
    t= t + s.charAt(k) + s.charAt(k);

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k= k+1;  
// if s[k-1] has a twin in s, add 1 to k  
if (k < s.length() && s.charAt(k-1) == s.charAt(k)) {  
    k= k+1;  
}
```